CS5112: Algorithms and Data Structures for Applications

Lecture 14: Exponential decay; convolution

Ramin Zabih

Some content from: Piotr Indyk; Wikipedia/Google image search;

J. Leskovec, A. Rajaraman, J. Ullman: Mining of Massive Datasets, http://www.mmds.org





Administrivia

- Q7 delayed due to Columbus day holiday
- HW3 out but short, can do HW in groups of 3 from now on
- Class 10/23 will be prelim review for 10/25
 - Greg lectures on dynamic programming next week
- Anonymous survey out, please respond!
- Automatic grading apology



Survey feedback





Automatic grading apology

- In general students should not have a grade revised downward
 - Main exception is regrade requests
- Automatic grading means that this sometimes happened
- At the end, we decided that the priority was to assign HW grades based on how correct the code was
- Going forward, please treat HW grades as tentative grades
 - We will announce when those grades are finalized
 - After, they will only be changed under exceptional circumstances



Today

- Two streaming algorithms
- Convolution



Some natural histogram queries



How many elements have non-zero frequency?



Streaming algorithms (recap)

- 1. Boyer-Moore majority algorithm
- 2. Misra-Gries frequent items
- 3. Find popular recent items (box filter)
- 4. Find popular recent items (exponential window)
- 5. Flajolet-Martin number of items



Weighted average in a sliding window

- Computing the average of the last k inputs can be viewed as a dot product with a constant vector $v = \left[\frac{1}{\nu}, \frac{1}{\nu}, \dots, \frac{1}{\nu}\right]$
- Sometimes called a box filter
 - Easy to visualize
- This is also a natural way to smooth, e.g., a histogram

– To average together adjacent bins, $v = \left[\frac{1}{3}, \frac{1}{3}, \frac{1}{3}\right]$

• This kind of weighted average has a famous name: convolution



Decaying windows



- Let our input at time t be $\{a_1, a_2, \dots, a_t\}$
- With a box filter over all of these elements we computed



• Instead let us pick a small constant *c* and compute

$$\hat{a}_t = \sum_{i=0}^{t-1} a_{t-i} (1-c)^i$$



Easy to update this

• Update rule is simple, let the current dot product be \hat{a}_t

$$\hat{a}_{t+1} = (1-c)\hat{a}_t + a_{t+1}$$

- This downscales the previous elements correctly, and the new element is scaled by $(1 c)^0 = 1$
- This avoids falling off the edge
- Gives us an easy way to find popular items



4. Popular items with decaying windows

- We keep a small number of weighted sum counters
- When a new item arrives for which we already have a counter, we update it using decaying windows, and update all counters
- How do we avoid getting an unbounded number of counters?
- We set a threshold, say ½, and if any counter goes below that value we throw it away
- The number of counters is bounded by $\frac{2}{c}$



How many distinct items are there?

- This tells you the size of the histogram, among other things
- To solve this problem exactly requires space that is linear in the size of the input stream
 - Impractical for many applications
- Instead we will compute an efficient estimate via hashing



5. Flajolet-Martin algorithm

- Basic idea: the more different elements we see, the more different hash values we will see
 - We will pick a hash function that spreads out the input elements
 - Typically uses universal hashing



Flajolet-Martin algorithm

- Pick a hash function h that maps each of the n elements to at least log₂ n bits
- For input a, let r(a) be the number of trailing 0s in h(a)

-r(a) = position of first 1 counting from the right

-E.g., say h(a) = 12, then 12 is 1100 in binary, so r(a) = 2

- Record R = the maximum r(a) seen
- Estimated number of distinct elements = 2^R

- Anyone see the problem here?



Why It Works: Intuition

- Very rough intuition why Flajolet-Martin works:
 - -h(a) hashes a with equal probability to any of n values
 - Sequence of $(\log_2 n)$ bits; 2^{-r} fraction of a's have tail of r zeros
 - About 50% hash to ***0
 - About 25% hash to **00
 - So, if we saw the longest tail of r=2 (i.e., item hash ending *100) then we have probably seen about 4 distinct items so far
 - Hash about 2^r items before we see one with zero-suffix of length r



Convolution

- Weighted average with a mask/stencil/template
 - Dot product of vectors
- Many important properties and applications
- Symmetric in the inputs
- Equivalent to linear shift-invariant systems

 "Well behaved", in a certain precise sense
- Primary uses are smoothing and matching
- This is the "C" in "CNN"



Local averaging in action





Smoothing parameter effects





Matched filters

- Convolution can be used to find pulses
 - This is actually closely related to smoothing
- How do we find a known pulse in a signal? Convolve the signal with our template!
 - E.g. to find something in the signal that looks like [1 6 -10] we convolve with [1 6 -10]
- Question: what sense does this make?
 - Anecdotally it worked for finding boxes



Box finding example





Pulse finding example





Why does this work?

- Some nice optimality properties, but the way I described it, the algorithm fails
- Idea: the [1 6 -10] template gives biggest response when signal is [... 1 6 -10 ...]
 - Value is 137 at this point
- But is this actually correct?
 - You actually need both the template and the input to have a zero mean and unit energy (sum of squares)
 - Easily accomplished: subtract -1, then divide by 137, get 1/137 * [2 7 -9]



Geometric intuition

- Taking the dot product of two vectors
 - $-\operatorname{Recall} [a \ b \ c] \cdot [e \ f \ g] = ae \ + \ bf \ + \ cg$
 - Basically the projection of a vector on another



• The normalized vector with the biggest projection on x is, of course: x!

