

# Cryptocurrency Intro

# Cryptocurrency: what is it?

Fundamentally, it's a ledger.

Alice owes Bob \$10

Alice owes Charlie \$20

Charlie owes Bob \$50

Bob owes Deborah \$90

Deborah owes Alice \$15

# Cryptocurrency: what is it?

It's public.

Issues that arise:

- Privacy?
- How do we know both parties agree to a transaction? And how do we know people aren't impersonating others?
  - Digital Signatures -- more on this later!
- How do we know everyone will pay their debts?
  - Don't allow people to go net negative.

# Cryptocurrency: what is it?

It's public.

Alice has \$100

Bob has \$120

Charlie has \$50

Deborah has \$10

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Alice owes Bob \$10

Alice owes Charlie \$20

Charlie owes Bob \$50

Bob owes Deborah \$90

Deborah owes Alice \$15

# Cryptocurrency: what is it?

It's not tied to any other currency.

## CornellCoin Ledger

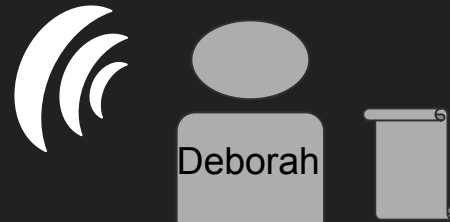
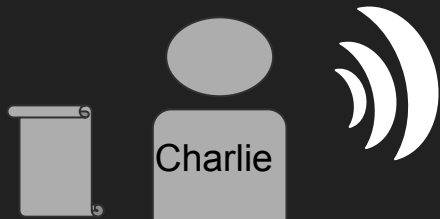
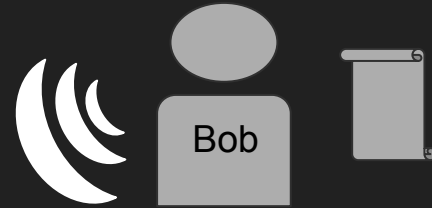
Alice has 100CC  
Bob has 120CC  
Charlie has 50CC  
Deborah has 10CC

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Alice owes Bob 10CC  
Alice owes Charlie 20CC  
Charlie owes Bob 50CC  
Bob owes Deborah 90CC

# Cryptocurrency: what is it?

It's decentralized.



# Cryptocurrency: what is it?

It's decentralized.

Issues that arise:

- What is the source of truth? How do we know that everyone has the same ledger?
  - Need to guarantee some kind of consensus.
- How is new money introduced to the system? And when? And to who?
  - “Work” -- more on this later!

# Cryptocurrency: what is it?

It's a ledger.

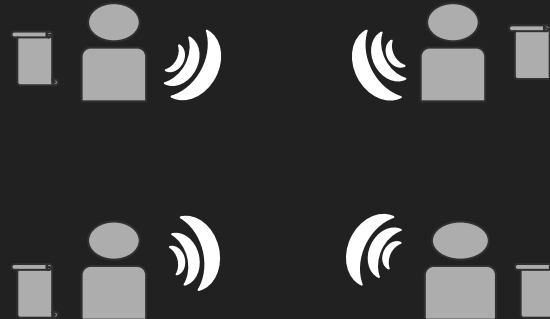
It's public.

It's decentralized.

CornellCoin Ledger	
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# Digital Signatures

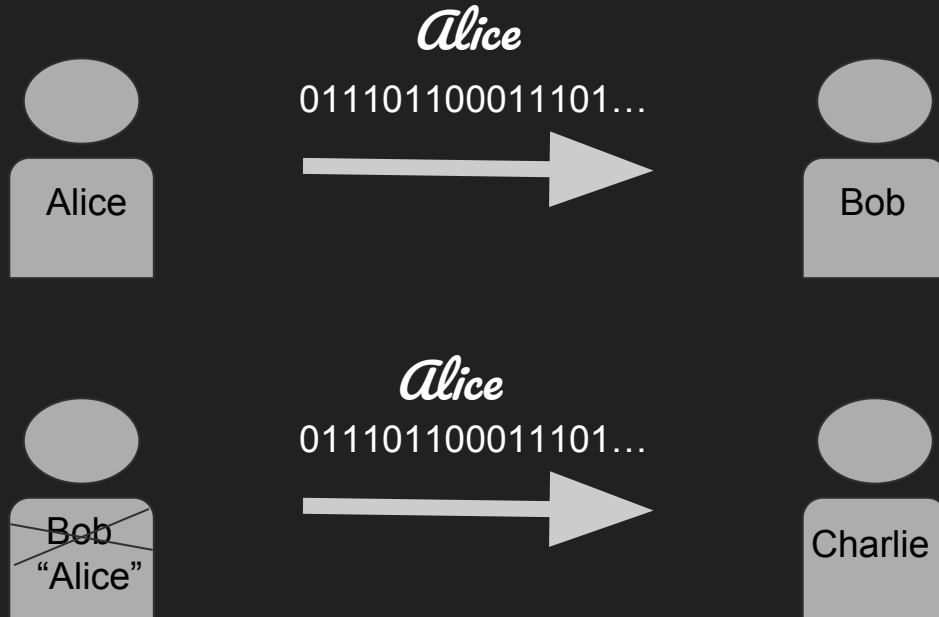
# Digital Signatures

Goal:

- Identify a particular person digitally
- Verifiable by a third party
- Not forgeable

# Digital Signatures

Why is this hard?



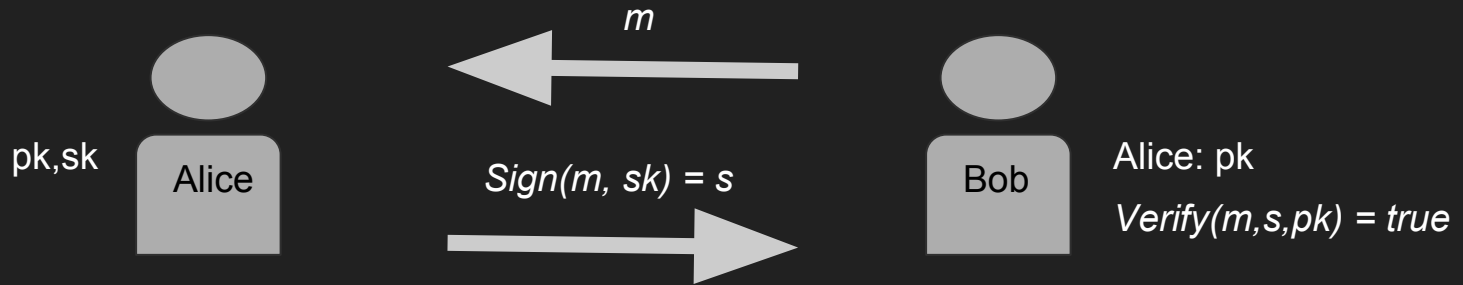
# Digital Signatures

1. Key Generator
  - a. Produces a Public Key and Private/Secret Key
  - b. Example: RSA
2.  $\text{Sign}(\text{message}, \text{private\_key}) \rightarrow \text{signature}$ 
  - a. Not reversible (without private\_key)
  - b. Output should appear uncorrelated with input
3.  $\text{Verify}(\text{message}, \text{signature}, \text{public\_key}) \rightarrow \text{boolean}$ 
  - a. Note: doesn't involve the private key!

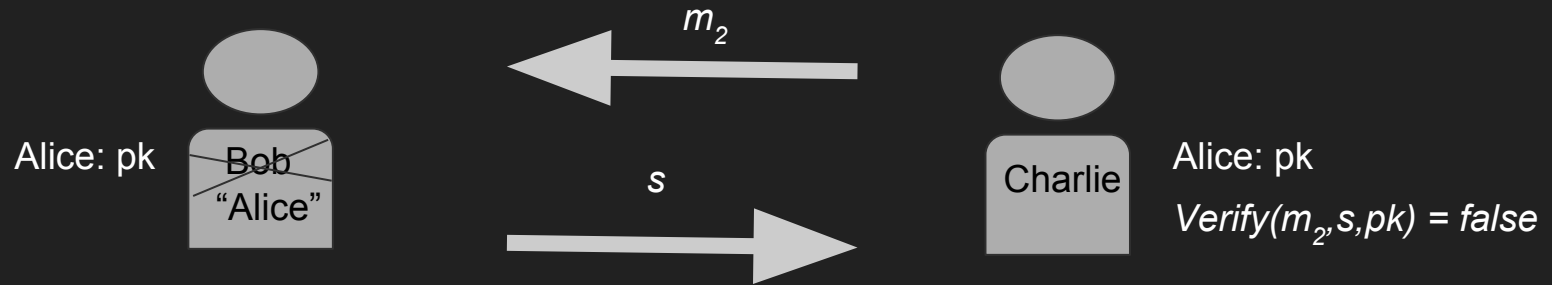
# Digital Signatures



# Digital Signatures



# Digital Signatures



# Digital Signatures

CornellCoin Ledger

Alice owes Bob 100CC    *Signed by Alice*



# Digital Signatures

## CornellCoin Ledger

Alice owes Bob 100CC      *Signed by Alice*

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Alice owes Bob 100CC      *Signed by Alice*

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Alice owes Bob 100CC      *Signed by Alice*

# Digital Signatures

## CornellCoin Ledger

1. Alice owes Bob 100CC *Signed by Alice*
2. Alice owes Bob 100CC ~~*Signed by Alice*~~
3. Alice owes Bob 100CC ~~*Signed by Alice*~~
4. Alice owes Bob 100CC ~~*Signed by Alice*~~
5. Alice owes Bob 100CC ~~*Signed by Alice*~~

# Digital Signatures: Implementation

- Implementation is for fun!
  - Presenting a (simplified) version of the RSA implementation
- Prime numbers
- Modulo arithmetic
  - Reminder: this is getting the remainder after division
  - $10 \bmod 4 = 2$
  - $15 \bmod 5 = 0$
- Goal: Find  $e, d, n$ , such that  $(m^e)^d \equiv m \pmod{n}$ 
  - Should also be tricky to determine  $d$  given  $e, n, m$
- Examples taken from Wikipedia

OPTIONAL MATERIAL

# Digital Signatures: KeyGen Implementation

- Choose two prime numbers:  $p, q$ 
  - Should be chosen randomly
- Compute  $n = pq$
- Compute  $\lambda(n) = LCM((p-1, q-1))$
- Choose  $1 < e < \lambda(n)$ , such that  $e$  and  $\lambda(n)$  are coprime
- Solve for  $d$ :  $d * e = 1 \pmod{\lambda(n)}$
  
- Public key:  $(n, e)$
- Private key:  $(n, d)$

OPTIONAL MATERIAL

# Digital Signatures: Implementation

- Reminder: Find  $e, d, n$ , such that  $(m^e)^d \equiv m \pmod{n}$
- Public key:  $(n, e)$     Private key:  $(n, d)$
  
- *Sign* $(m, sk): m^d \pmod{n}$
- *Verify* $(m, s, pk): s^e \pmod{n} == m$

OPTIONAL MATERIAL

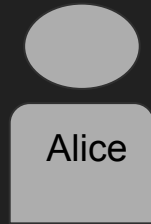
# Digital Signatures: KeyGen Implementation Example

- $p = 61, q = 53$
- $n = pq = 61 * 53 = 3233$
- $\lambda(3233) = LCM(60, 52) = 780$
- Choose  $e = 17$
- Solve for  $d$ :  $d * 17 = 1 \text{ mod } 3233 \rightarrow d = 413$
  
- Public key:  $(3233, 17)$
- Private key:  $(3233, 413)$

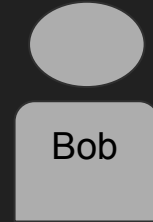
OPTIONAL MATERIAL

# Digital Signatures: Example

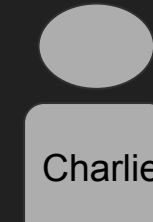
pk = (3233, 17)  
sk = (3233, 413)



pk: (3233, 17)



Alice: (3233, 17)



Alice: (3233, 17)

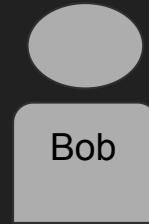
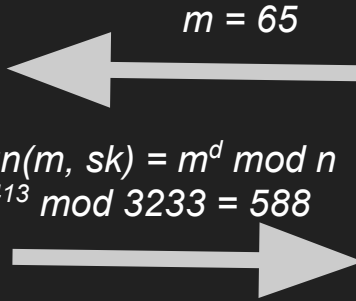
OPTIONAL MATERIAL

# Digital Signatures: Example

pk = (3233, 17)  
sk = (3233, 413)



$Sign(m, sk) = m^d \bmod n$   
 $65^{413} \bmod 3233 = 588$



Alice: (3233, 17)

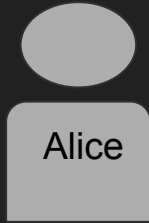
$Verify(m, s, pk) = (s^e \bmod n == m)$   
 $588^{17} \bmod 3233 == 65$

OPTIONAL MATERIAL



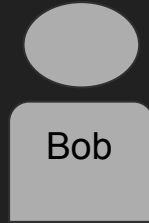
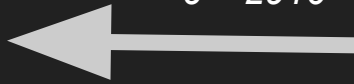
# Public/Private Key Encryption

pk = (3233, 17)  
sk = (3233, 413)



Alice

$c = 2910$



Bob

Alice: (3233, 17)

$m = 92$

$Decrypt(c, sk) = c^d \bmod n$   
 $2910^{413} \bmod 3233 = 92$

$Encrypt(m, pk) = m^e \bmod n$   
 $92^{17} \bmod 3233 = 2910$

OPTIONAL MATERIAL

# Digital Signatures

Unanswered questions:

- Why on earth does this work?
  -
- How do Bob and Charlie know that the public key they received is actually from Alice?
  - Practically... not a technical solution. Certificate Authorities do the job.
  - Can the blockchain be used for this?